## Tradeoffs between Power Management and Tail Latency in Warehouse-Scale Applications



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## Warehouse-scale computers (WSC)

Datacenters built for a specific class of workloads

Heterogeneous, multi-tiered distributed services, tightly coupled

Overall service must provide latency guarantees often in the ~100ms



[Meisner et al. 2011]

# WSC performance metric is most often tail latency

	50%ile latency	95%ile latency	99%ile latency
One random leaf finishes	lms	5ms	10ms
95% of all leaf requests finish	12ms	32ms	70ms
100% of all leaf requests finish	40ms	87ms	140ms

[Dean et al. 2012]

#### Many services require a response from all leaves

Which could be orders of magnitude slower than average responses

And very sensitive to variability

## Power management leads to performance variability



A space between power savings and worse tail latency

### **Opportunities for power management**



### Energy proportionality



Energy proportionality: scale server power with load

Stable across platform generations

## Energy proportionality



Energy proportionality: scale server power with load

Worst in mid-range utilization

#### Servers see the full range of utilization



[Content ads cluster in North America]

15-100% swings during a single day



idle (W)

full load (W)

Processors are still the major power consumers but cores also scale best with load

At low-utilization, non-proportional components (disks, flash, DRAM) matter more

#### Idle power management



## Processor idle power management (C-states)

OS-exposed mechanism to exploit idle periods but still HW-controlled

Mostly responsible for current processors' proportionality

Various degrees of power gating increasing power savings increasing wakeup latency [1-200 µs]



#### Some WSC workloads sleep in short bursts



#### Some WSC workloads sleep in short bursts



Application sleep activity length can be comparable to wakeup latencies  $\rightarrow$  deep sleep can hurt service latency

#### Effects of sleep state selection



Deep sleep does save significant power (up to 15%)

But also hurts tail latency (up to 15%)

#### Active power management



#### WSC services are often stalled on memory



A good candidate for voltage and frequency scaling (DVFS)

## Wishlist for server DVFS

#### Zero tolerance

latency degradation is evil

#### Workload independence

thousands of relevant workloads

Fine-grained requests handled in O(1ms)

#### Per-core

scalable services have independent threads handling independent requests

#### Importance of fine granularity



No power savings for control as fast as 100 µs

Execution phases are likely more fine-grained and would be best exploited in hardware

#### Importance of workload (in)dependence



## Takeaways



Current server hardware is not universally energy proportional. Especially related to components like flash, DRAM, or disks.



Core sleep states (clock & power gating) are mostly responsible for power savings. But their effects on latency should be treated with care.



Active power savings are possible either on a very fine granularity, with additional hardware, or on ubiquitous individual workloads, exploiting latency slack.

#### Servers are often underutilized



#### Operating in power inefficient regions

[Barroso et al., 2013, several thousand machines over 3 months]

#### ... but also require a lot of computation



[Reddi et al. 2010, Bing websearch]

#### Some services can be overdesigned



Specifically, to handle the peak utilization case

There is no benefit in beating service agreements (SLAs) at low utilization

#### Energy proportionality



Energy proportionality: scale server power with load

Relatively stable across platform generations